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The 6-girth-thickness of the complete graph

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ABSTRACT

The g -girth-thickness $\theta(g, G)$ of a graph G is the minimum number of planar subgraphs of girth at least g whose union is G . In this paper, we determine the 6-girth-thickness $\theta(6, K_n)$ of the complete graph K_n in almost all cases. And also, we calculate by computer the missing value of $\theta(4, K_n)$.

KEYWORDS

Thickness; planar decomposition; complete graph; girth

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1. Introduction

In this paper, all graphs are finite and simple. A graph in which any two vertices are adjacent is called a *complete graph* and it is denoted by K_n if it has n vertices. If a graph can be drawn in the Euclidean plane such that no inner point of its edges is a vertex or lies on another edge, then the graph G is called *planar*. The *girth* of a graph is the size of its shortest cycle or ∞ if it is acyclic. It is known that an acyclic graph of order n has size at most $n - 1$ and a planar graph of order n and finite girth g has size at most $\frac{g-2}{g-2}(n - 2)$, see [8].

The *thickness* $\theta(G)$ of a graph G is the minimum number of planar subgraphs whose union is G . Equivalently, it is the minimum number of colors used in any edge coloring of G such that each set of edges in the same chromatic class induces a planar subgraph.

The concept of the thickness was introduced by Tutte [19]. The problem to determine the thickness of a graph G is NP-hard [15], and only a few of exact results are known, for instance, when G is a complete graph [2, 5, 6], a complete multipartite graph [7, 11, 18, 21, 22] or a hypercube [14].

Generalizations of the thickness for the complete graphs also have been studied such that the *outerthickness* θ_o , defined similarly but with outerplanar instead of planar [12], and the *S-thickness* θ_S , considering the thickness on a surface S instead of the plane [4]. The thickness has many applications, for example, in the design of circuits [1], in the

Ringel's earth-moon problem [13], or to bound the achromatic numbers of planar graphs [3]. See also [16].

In [17], the *g-girth-thickness* $\theta(g, G)$ of a graph G was defined as the minimum number of planar subgraphs of girth at least g whose union is G . Indeed, the g -girth thickness generalizes the thickness when $g = 3$ and the *arboricity number* when $g = \infty$.

This paper is organized as follows. In Section 2, we obtain the 6-girth-thickness $\theta(6, K_n)$ of the complete graph K_n getting that $\theta(6, K_n)$ equals $\lceil \frac{n+2}{3} \rceil$, except for $n = 3t + 1, t \geq 4$ and $n \neq 2$, for which $\theta(6, K_2) = 1$. In Section 3, we show that there exists a set of 3 planar triangle-free subgraphs of K_{10} whose union is K_{10} . The decomposition was found by computer and, as a consequence, we disproved the conjecture that appears in [17] about the missing case of the 4-girth-thickness of the complete graph.

2. Determining $\theta(6, K_n)$

A planar graph of n vertices with girth at least 6 has size at most $3(n - 2)/2$ for $n \geq 6$ and size at most $n - 1$ for $1 \leq n \leq 5$, therefore, the 6-girth-thickness $\theta(6, K_n)$ of the complete graph K_n is at least

$$\left\lceil \frac{n(n-1)}{3(n-2)} \right\rceil = \left\lceil \frac{n+1}{3} + \frac{2}{3n-6} \right\rceil = \left\lceil \frac{n+2}{3} \right\rceil$$

for $n \geq 6$, as well as, $\lceil \frac{n+2}{3} \rceil$ for $n \in \{1, 3, 4, 5\}$. We have the following theorem.

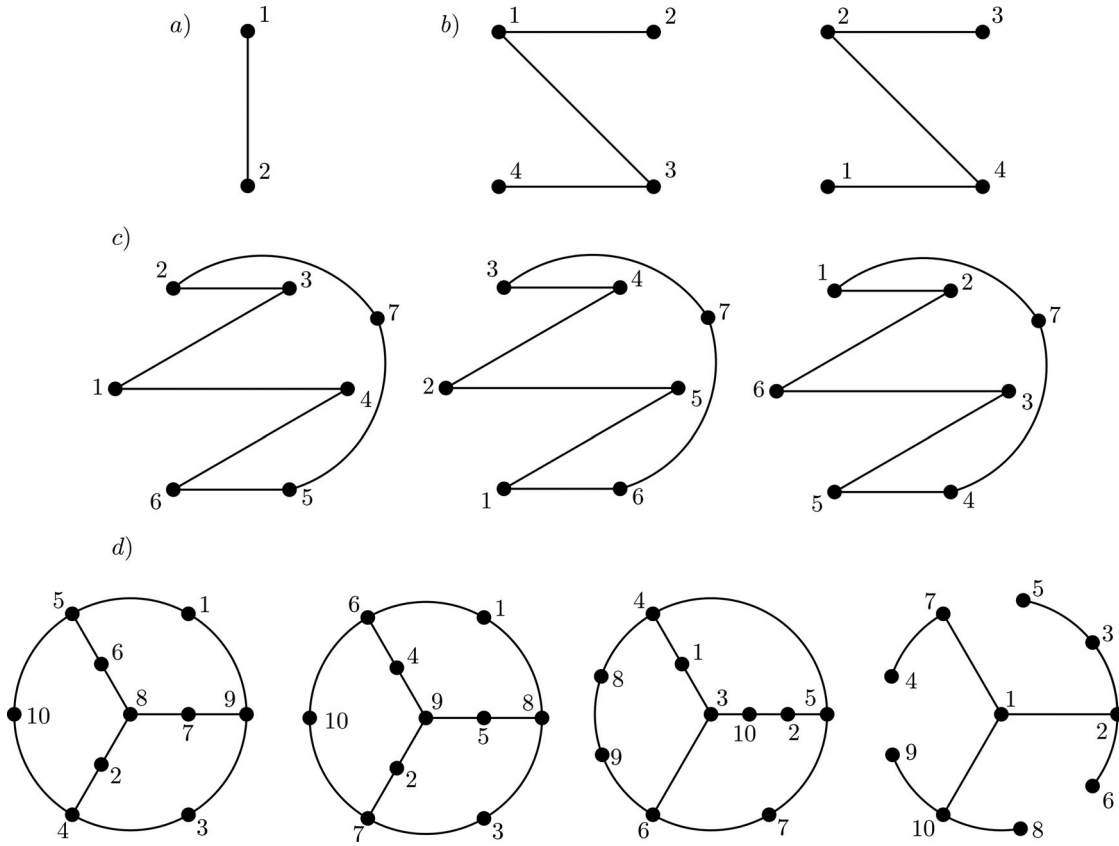


Figure 1. A decomposition of K_n into $\theta(6, K_n)$ planar subgraphs of girth at least 6: (a) for $n=2$, (b) for $n=4$, (c) for $n=7$ and (d) for $n=10$.

Theorem 2.1. *The 6-girth-thickness $\theta(6, K_n)$ of K_n is equal to $\lceil \frac{n+2}{3} \rceil$ except possibly when $n = 3t + 1$, for $t \geq 4$, and $n \neq 2$ for which $\theta(6, K_2) = 1$.*

Proof. To begin with, Figure 1 displays equality for $n = 2, 4, 7, 10$ with $\theta(6, K_n) = 1, 2, 3, 4$, respectively. The rest of the cases for $1 \leq n \leq 10$ are obtained by the hereditary property of the induced subgraphs. We remark that the decomposition of K_{10} was found by computer using the database of the connected planar graphs of order 10 that appears in [9].

Now, we need to distinguish two main cases, namely, when t is even or t is odd for $n = 3t$, that is, when $n = 6k$ and $n = 6k + 3$ for $k \geq 2$. The cases $n = 6k - 1$ and $n = 6k + 2$, i.e., for $n = 3t + 1$, are obtained by the hereditary property of the induced subgraphs, that is, since $K_{6k-1} \subset K_{6k}$ and $K_{6k+2} \subset K_{6k+3}$, we have

$$2k + 1 \leq \theta(6, K_{6k-1}) \leq \theta(6, K_{6k}) \text{ and} \\ 2k + 2 \leq \theta(6, K_{6k+2}) \leq \theta(6, K_{6k+3}), \text{ respectively.}$$

Therefore, the case of $n = 6k$ shows a decomposition of K_{6k} into $2k + 1$ planar subgraphs of girth at least 6, while the case of $n = 6k + 3$ shows a decomposition of K_{6k+3} into $2k + 2$ planar subgraphs of girth at least 6. Both constructions are based on the planar decomposition of K_{6k} of Beineke and Harary [5] (see also [2, 6, 20]) but we use the combinatorial approach given in [3]. Then, for the sake of completeness, we give a decomposition of K_{6k} in order to obtain its usual thickness. In the remainder of this proof, all sums are taken modulo $2k$.

We recall that complete graphs of even order $2k$ are decomposable into a cyclic factorization of Hamiltonian paths, see [10]. Let G^x be a complete graph of order $2k$, label its vertex set $V(G^x)$ as $\{x_1, x_2, \dots, x_{2k}\}$ and let \mathcal{F}_i^x be the Hamiltonian path with edges

$$x_i x_{i+1}, x_{i+1} x_{i-1}, x_{i-1} x_{i+2}, x_{i+2} x_{i-2}, \dots, x_{i+k+1} x_{i+k},$$

for all $i \in \{1, 2, \dots, k\}$. The partition $\{E(\mathcal{F}_1^x), E(\mathcal{F}_2^x), \dots, E(\mathcal{F}_k^x)\}$ is such factorization of G^x . We remark that the center of \mathcal{F}_i^x has the edge $e_i^x = x_{i+\lceil \frac{k}{2} \rceil} x_{i+\lfloor \frac{k}{2} \rfloor}$, see Figure 2.

Let G^u, G^v and G^w be the complete subgraphs of K_{6k} having $2k$ vertices each of them and such that G^w is $K_{6k} \setminus (V(G^u) \cup V(G^v))$. The vertices of $V(G^u), V(G^v)$ and $V(G^w)$ are labeled as $\{u_1, u_2, \dots, u_{2k}\}, \{v_1, v_2, \dots, v_{2k}\}$ and $\{w_1, w_2, \dots, w_{2k}\}$, respectively.

Let x be an element of $\{u, v, w\}$. Take the cyclic factorization $\{E(\mathcal{F}_1^x), E(\mathcal{F}_2^x), \dots, E(\mathcal{F}_k^x)\}$ of G^x into Hamiltonian paths and denote as P_{x_i} and $P_{x_{i+k}}$ the subpaths of \mathcal{F}_i^x containing k vertices and the leaves x_i and x_{i+k} , respectively. We define the other leaves of P_{x_i} and $P_{x_{i+k}}$ as $f(x_i)$ and $f(x_{i+k})$, respectively and according to the parity of k , that is (see Figure 2),

$$f(x_i) = \begin{cases} x_{i+\lceil \frac{3k}{2} \rceil} & \text{if } k \text{ is odd,} \\ x_{i+\lceil \frac{k}{2} \rceil} & \text{if } k \text{ is even.} \end{cases} \text{ and} \\ f(x_{i+k}) = \begin{cases} x_{i+\lceil \frac{k}{2} \rceil} & \text{if } k \text{ is odd,} \\ x_{i+\lceil \frac{3k}{2} \rceil} & \text{if } k \text{ is even.} \end{cases}$$

We remark that the set of edges $\{x_i x_{i+k} : 1 \leq i \leq k\}$ is the same set of edges that $\{f(x_i) f(x_{i+k}) : 1 \leq i \leq k\}$.

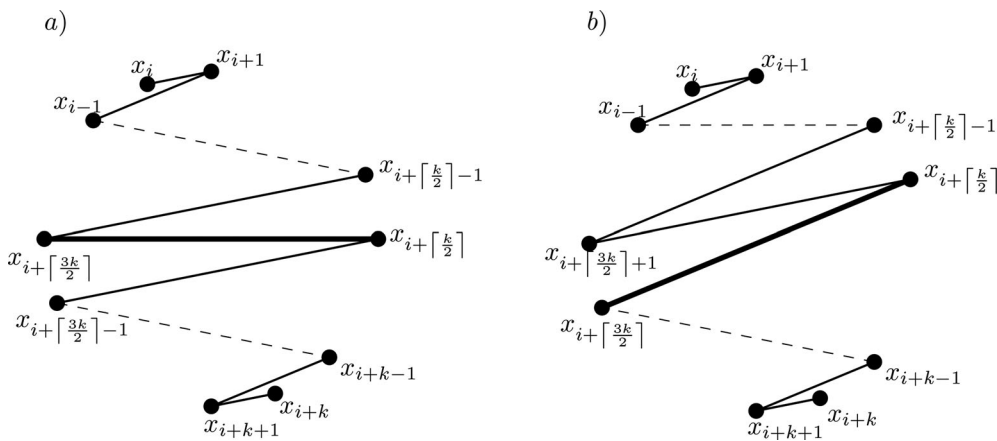


Figure 2. The Hamiltonian path \mathcal{F}_i^x : Left (a) The edge e_i^x in bold for k odd. Right (b) The edge e_i^x in bold for k even.

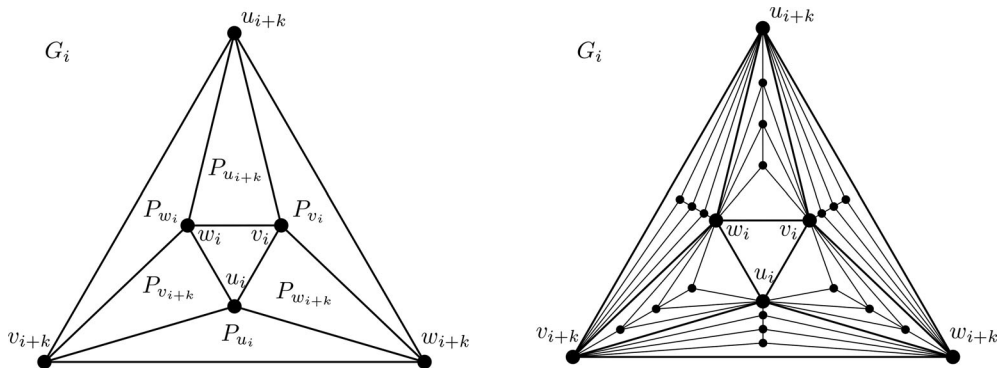


Figure 3. (Left) The octahedron subgraph of the graph G_i . (Right) The graph G_i .

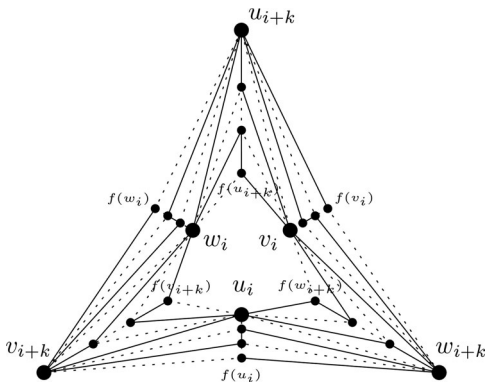


Figure 4. Partial modification of the subgraph G_i .

Now, we construct the maximal planar subgraphs G_1, G_2, \dots, G_k and a matching G_{k+1} with $6k$ vertices each in the following way. Let G_{k+1} be the perfect matching with the edges $u_j u_{j+k}, v_j v_{j+k}$ and $w_j w_{j+k}$ for $j \in \{1, 2, \dots, k\}$.

For each $i \in \{1, 2, \dots, k\}$, let G_i be the spanning planar graph of K_{6k} whose adjacencies are given as follows: we take the 6 paths, $P_{u_i}, P_{u_{i+k}}, P_{v_i}, P_{v_{i+k}}, P_{w_i}$ and $P_{w_{i+k}}$ and insert them in the octahedron with the vertices $u_i, u_{i+k}, v_i, v_{i+k}, w_i$ and w_{i+k} as is shown in Figure 2 (Left). The vertex x_j of each path P_{x_j} is identified with the vertex x_j in the corresponding triangle face and join all the other vertices of the path with both of the other vertices of the triangle face, see Figure 3 (Right).

By construction of $G_i, K_{6k} = \cup_{i=1}^{k+1} G_i$, see [2, 5] to check a full proof. In consequence, the $k+1$ planar subgraphs G_i show that $\theta(3, K_{6k}) \leq k+1$ and then, $\theta(3, K_{6k}) = k+1$ owing to the fact that $\theta(3, K_{6k}) \geq \left\lceil \frac{\binom{6k}{2}}{3(6k-2)} \right\rceil = k+1$.

Now, we proceed to prove that $\theta(6, K_{6k}) \leq 2k+1$ in Case 1 and $\theta(6, K_{6k+3}) \leq 2k+2$ in Case 2. The main idea of both cases is divide each G_i into two subgraphs of girth 6 for any $i \in \{1, \dots, k\}$.

- Case $n = 6k$. Consider the set of planar subgraphs $\{G_1, G_2, \dots, G_{k+1}\}$ of K_{6k} which is described above.

Step 1. For each $i \in \{1, \dots, k\}$, remove the six edges of the triangles $u_i v_i w_i$ and $u_{i+k} v_{i+k} w_{i+k}$.

Step 2. For each $i \in \{1, \dots, k\}$, divide the obtained subgraph into two subgraphs H_i^1 and H_i^2 as follows: The maximum matching of P_{x_i} incident to the vertex $f(x_i)$ belongs to H_i^1 (see dotted subgraph in Figure 4) while the maximum matching of $P_{x_{i+k}}$ incident to the vertex $f(x_{i+k})$ belongs to H_i^2 .

Next, the rest of the edges joined to the vertices of the paths P_{x_i} and $P_{x_{i+k}}$, in an alternative way from the exterior region to the region with the vertices $\{u_i, v_i, w_i\}$, belong to H_i^1 and H_i^2 respectively, such that the edges $f(w_i)u_{i+k}, f(v_i)w_{i+k}$ and $f(u_i)v_{i+k}$ belong to H_i^1 and the edges $f(w_i)v_{i+k}, f(v_i)u_{i+k}$ and $f(u_i)w_{i+k}$ belong to H_i^2 , see Figure 4.

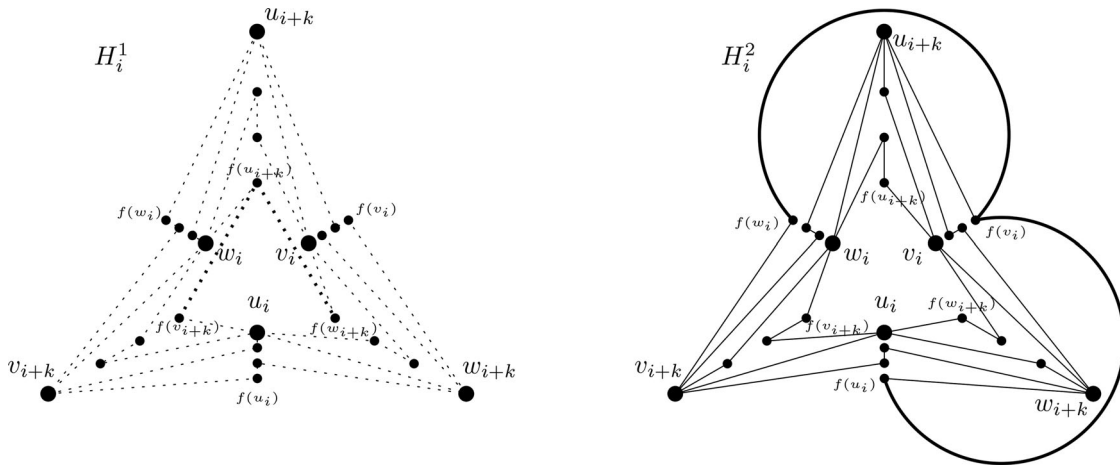


Figure 5. Subgraphs H_i^1 and H_i^2 for the Case 1.

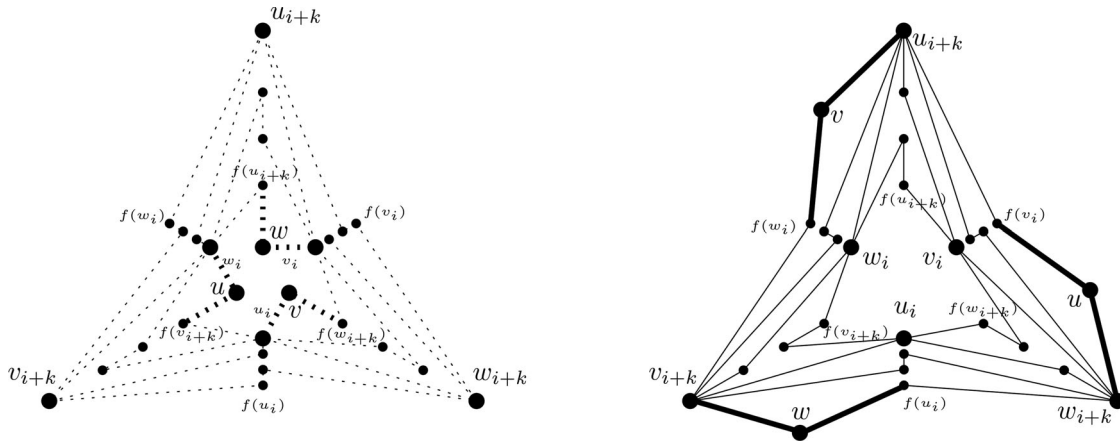


Figure 6. Subgraphs H_i^1 and H_i^2 for Case 2.

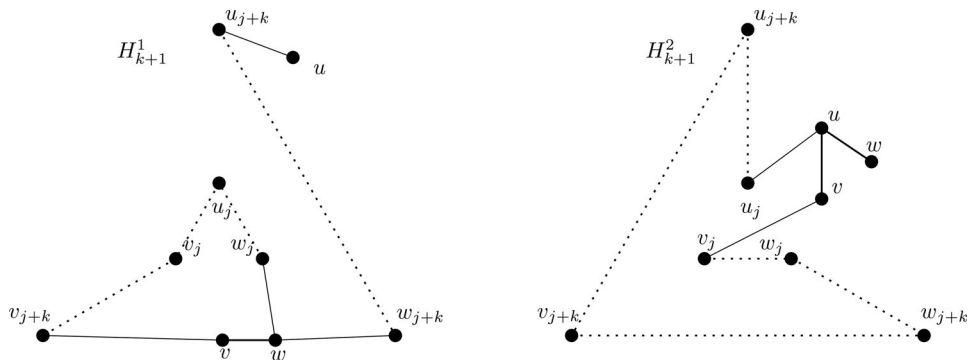


Figure 7. Partial subgraphs H_{k+1}^1 and H_{k+1}^2 .

Step 3. Consider the removed edges in Step 1, add the edges $f(v_{i+k})f(u_{i+k})$ and $f(u_{i+k})f(w_{i+k})$ to H_i^1 and the edges $f(w_i)f(v_i)$ and $f(v_i)f(u_i)$ to H_i^2 , see Figure 5. The rest of the edges removed in Step 1 are added to G_{k+1} getting the subgraph H_{k+1} which is the union of the paths $\{f(v_i), f(v_{i+k}), f(w_{i+k}), f(w_i), f(u_i), f(u_{i+k})\}$.

2. Case $n = 6k + 3$. Consider the set of planar subgraphs $\{G_1, G_2, \dots, G_{k+1}\}$ of K_{6k} which is described above as well as Step 1 and 2 of the previous case.

Step 3. Add three vertices u, v and w in the subgraphs H_i^1 and H_i^2 , for each $i \in \{1, \dots, k\}$, and the edges $uw, uv, vw, uv, vw, uv, vw$ into H_i^1 as well as the edges $uw, uv, vw, uv, vw, uv, vw$ into H_i^2 , see Figure 6.

Step 4. On one hand, remains to define the adjacencies between u, v, w and all the adjacencies between u and u_i, v and v_i, w and w_i for each $j \in \{1, \dots, k\}$. On the other hand, the edges of the graph G_{k+1} together with the removed edges of the Step 1 form a set of triangle

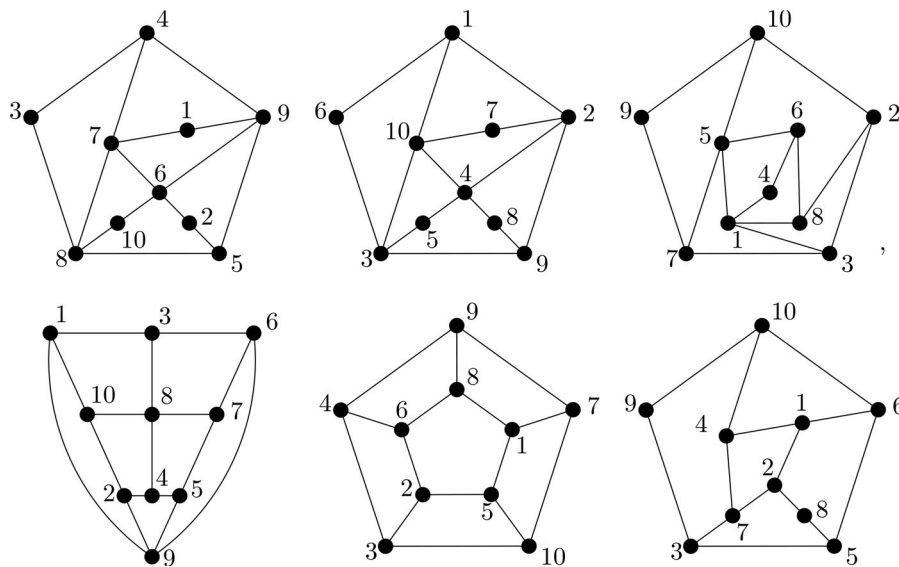


Figure 8. Two planar decompositions of K_{10} into three subgraphs of girth 4.

prisms which we split into two subgraphs called H_{k+1}^1 and H_{k+1}^2 in the following way:

- The adjacency vw is in H_{k+1}^1 while the adjacencies uv and uw are in H_{k+1}^2 , see Figure 7.
- The set of adjacencies vv_{j+k} , ww_j , ww_{j+k} and uu_{j+k} are in H_{k+1}^1 while the set of adjacencies vv_j and uu_j are in H_{k+1}^2 , for each $j \in \{1, \dots, k\}$, see Figure 7.
- The subgraph H_{k+1}^1 contains the adjacencies $v_{j+k}v_j$, v_ju_j , u_jw_j and $w_{j+k}u_{j+k}$ (a set of subgraphs $P_4 \cup K_2$) and the subgraph H_{k+1}^2 contains the adjacencies u_ju_{j+k} , $u_{j+k}v_{j+k}$, $v_{j+k}w_{j+k}$, $w_{j+k}w_j$ and w_jv_j (a set of subgraphs P_6) for all $j \in \{1, \dots, k\}$, see Figure 7.

By the small cases and the two main cases, the theorem follows. \square

3. The 4-girth thickness of K_{10}

In [17], Rubio-Montiel gave a decomposition of K_n into $\theta(4, K_n) = \lceil \frac{n+2}{4} \rceil$ triangle-free planar subgraphs, except for $n=10$. In that case, it was bounded by $3 \leq \theta(4, K_{10}) \leq 4$ and conjectured that the correct value was the upper bound. Using the database of the connected planar graphs of order 10 that appears in [9] and the SageMath program, we found two decompositions of K_{10} into 3 planar subgraphs of girth at least 4 illustrated in Figure 8. In summary, the correct value of $\theta(4, K_n)$ was the lower bound and then, we have the following theorem.

Theorem 3.1. *The 4-girth-thickness $\theta(4, K_n)$ of K_n equals $\lceil \frac{n+2}{4} \rceil$ for $n \neq 6$ and $\theta(4, K_6) = 3$.*

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Disclosure statement

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