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# Hamiltonian claw-free graphs with locally disconnected vertices

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# ABSTRACT

An edge of *G* is *singular* if it does not lie on any triangle of *G*; otherwise, it is *non-singular*. A vertex *u* of a graph *G* is called *locally connected* if the induced subgraph G[N(u)] by its neighborhood is connected; otherwise, it is called *locally disconnected*.

In this paper, we prove that if a connected claw-free graph *G* of order at least three satisfies the following two conditions: For each locally disconnected vertex v of *G* with degree at least 3, there is a nonnegative integer *s* such that v lies on an induced cycle of length at least 4 with at most *s* non-singular edges and with at least s - 3 locally connected vertex; for each locally disconnected vertex v of *G* with degree 2, there is a nonnegative integer *s* such that v lies on an induced cycle *C* with at most *s* non-singular edges and with at least s - 2 locally connected vertices and such that the subgraph induced by those vertices of *C* that have degree two in *G* is a path or a cycle, then *G* is Hamiltonian, and it is best possible in some sense.

Our result is a common extension of two known results in Bielak (2000) and in Li (2002) ; hence also of the results in Oberly and Sumner (1979) and in Ryjáček (1990).

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# 1. Introduction

We consider only finite undirected simple graphs, unless otherwise stated. For terminology and notation not defined in this paper we refer to [9].

If *H* is a graph, then the *line graph* of *H*, denoted by L(H), is the graph with E(H) as its vertex set, in which two vertices are adjacent if and only if the corresponding edges have a vertex in common. For a family  $\mathcal{F}$  of a connected graphs, a graph is called  $\mathcal{F}$ -free if it contains no induced copies of any member of  $\mathcal{F}$ . The graph  $K_{1,3}$  is called a *claw*. It is a well-known fact that every line graph is claw-free, hence the class of the claw-free graphs can be considered as a natural generalization of the class of line graphs.

The neighborhood of a vertex v in G is denoted by  $N_G(v)$ . Denote  $N_G[v] = N_G(v) \cup \{v\}$ . A vertex v of G is locally connected if  $G[N_G(v)]$  is connected; otherwise, it is locally disconnected. Let LC(G) denote the set of all locally connected vertices of G. A graph G is called *locally connected* if every vertex of G is locally connected, *i.e.*, LC(G) = V(G). Oberly and Sumner proved the following well-known result.

**Theorem 1** (Oberly and Sumner [5]). Every connected, locally connected claw-free graph on at least three vertices is Hamiltonian.

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**Fig. 1.** The graphs  $G_1$ ,  $G_2$ ,  $G_3$  and  $G_4$ .

We say that a vertex v of a graph G is  $N_2$ -locally connected if the subgraph of G induced by the edge set  $\{e = xy \in E(G) : v \notin \{x, y\}$  and  $\{x, y\} \cap N(v) \neq \emptyset\}$  is connected. A graph G is called  $N_2$ -locally connected if every vertex of G is  $N_2$ -locally connected. It follows from the definitions that every locally connected graph is  $N_2$ -locally connected, but the converse is not true.

In 1990, Ryjáček [7] considered the graphs with some locally disconnected vertices in claw-free graphs and strengthened Theorem 1 by using this concept of  $N_2$ -locally connected. He showed that every connected  $N_2$ -locally connected claw-free graph *G* with  $\delta(G) \ge 2$  satisfying that *G* has no induced subgraph *H* isomorphic to either  $G_1$  or  $G_2$  (in Fig. 1) such that every vertex of degree 4 in *H* is locally disconnected in *G* is Hamiltonian. Bielak later improved this result by weakening the condition. Their result can be restated as the following theorem, where  $V_i(G) = \{x : d_G(x) = i\}$  and  $V_{\ge i}(G) = \{x : d_G(x) \ge i\}$ .

**Theorem 2** (Bielak [1]). Let G be a connected, N<sub>2</sub>-locally connected claw-free graph with  $\delta(G) \geq 2$  such that

(1) every induced subgraph H of G isomorphic to either  $G_1$  or  $G_2$  (in Fig. 1) has at least one locally connected vertex of G in  $V_3(H) \cup V_4(H)$ .

# Then G is Hamiltonian.

In this paper, we shall continue to extend the above result which will need some notation. We say that a vertex v of a graph G is  $N^2$ -locally connected if the subgraph of G induced by the vertices { $x \in V(G) : 1 \le d(x, v) \le 2$ } is connected, where d(x, v) denotes the distance between x and v. A graph G is called  $N^2$ -locally connected if every vertex of G is  $N^2$ -locally connected. Obviously, every  $N_2$ -locally connected graph is  $N^2$ -locally connected, but the converse is not generally true.

**Theorem 3.** Let *G* be a connected,  $N^2$ -locally connected claw-free graph with  $\delta(G) \ge 2$  satisfying

(2) every induced subgraph isomorphic to one of  $\{G_1, G_2, G_3, G_4\}$  (in Fig. 1) has at least one locally connected vertex of G in  $V_3(H) \cup V_4(H)$ .

Then G is Hamiltonian.

From Theorem 3, one can obtain the following known result immediately.

**Corollary 4** (*Li*[4]). Every connected  $N^2$ -locally connected { $G_1, G_2, G_3, G_4, K_{1,3}$ }-free graph G with  $\delta(G) \ge 2$  is Hamiltonian.

Let  $G_0$  be the graph obtained from some graph  $G_i$  in Fig. 1 by joining all vertices of an additional complete graph of arbitrarily larger order to some vertex of degree four or three in  $G_i$  and to its neighbors. Then  $G_0$  satisfies the conditions of Theorem 3 but not Corollary 4. This shows that Theorem 3 is stronger than Corollary 4.

Motivated by the above observation, in this paper, we intend to generality them by avoiding using the concept of  $N_2$ -(or  $N^2$ -)connected and use certain technical conditions on locally disconnected vertices instead. Here we need divide all edges of the graphs into two kinds of edges: An edge *e* of *G* is *singular* if it does not lie on any triangle of *G*; otherwise, it is *non-singular*. We have the following result that can deduce Theorem 3, as showed in Section 4.

#### Theorem 5. Let G be a connected claw-free graph of order at least three such that

- (i) for each locally disconnected vertex v of degree at least 3 in *G*, there is a nonnegative integer *s* such that v lies on an induced cycle of length at least four with at most *s* non-singular edges and with at least *s* 3 locally connected vertices;
- (ii) for each locally disconnected vertex v of degree 2 in G, there is a nonnegative integer s such that v lies on an induced cycle C with at most s non-singular edges and with at least s 2 locally connected vertices and such that  $G[V(C) \cap V_2(G)]$  is a path or a cycle.

# Then G is Hamiltonian.

In Section 2, we shall present Ryjáček's closure concept in claw-free graphs and some auxiliary results, which are then applied to the proof of our main result in Section 3. Section 4 is devoted to the proof of Theorem 3. In the last section, we discuss the sharpness of our main results, point out a flaw in the original proof of Corollary 4 and show that Theorem 5 is stronger than Theorem 3, and hence also than Corollary 4.

# 2. The closure of claw-free graphs

A locally connected vertex v is said to be *eligible* if  $G[N_G(v)]$  is not complete. For a vertex x of a graph G, the graph  $G_x^*$  with  $V(G_x^*) = V(G)$  and  $E(G_x^*) = E(G) \cup \{uv : u, v \in N_G(x)\}$  is called the *local completion* of G at x. For a claw-free graph G, let  $G_1 = G$ . For  $i \ge 1$ , if  $G_i$  is defined and if it has an eligible vertex  $x_i$ , then let  $G_{i+1} = (G_i)_{x_i}^*$ . If  $G_s = (G_{s-1})_{x_{s-1}}^*$  has no eligible vertex, then let  $cl(G) = G_s$  and call it the *closure of* G;  $G_1, \ldots, G_s$  is called a locally complete sequence of graphs that yields cl(G). The above operation was introduced in [8] and the following theorem sums up some properties.

**Theorem 6** (Ryjáček [8]). If G is a claw-free graph, then there is a closed claw-free graph cl(G) such that

- (3) the closure cl(*G*) is well-defined;
- (4) there is a triangle-free graph H such that cl(G) = L(H);
- (5) *G* is Hamiltonian if and only if cl(*G*) is Hamiltonian.

In a claw-free graph *G*, the locally disconnected vertices can be partitioned into three classes, depending on the structure of the graphs G[N(v)]: Let  $LD_i(G)$  be the set of locally disconnected vertices v for which there are exactly i components in G[N(v)] of order greater than one. (Note that the notations here are something different from [6].) Note that for a locally disconnected vertex v, G[N(v)] consists of exactly two complete subgraphs of *G*. Pfender proved the following.

**Lemma 7** (*Pfender*, [6]).  $(LD_0(cl(G)) \cup LD_1(cl(G))) \subseteq (LD_0(G) \cup LD_1(G))$  and  $LD_2(cl(G)) \subseteq LD_2(G)$  for every claw-freegraph G.

We need the following lemma, which follows from Lemma 7.

**Lemma 8.** For  $i \in \{0, 1, 2\}$ ,  $LD_i(cl(G)) \subseteq LD_i(G)$  for every claw-free graph G.

**Proof of Lemma 8.** Suppose that  $x \in LD_0(cl(G))$ , i.e, x is locally disconnected in cl(G) and  $d_{cl(G)}(x) = 2$ . By Lemma 7, x is locally disconnected in G, hence  $d_G(x) = 2$ . Thus  $x \in LD_0(G)$ .

Suppose that  $x \in LD_1(cl(G))$ . We claim that  $x \notin LD_0(G)$ : otherwise, let  $N_G(x) = \{y_1, y_2\}$ . Note that every edge in  $E(cl(G)) \setminus E(G)$  is non-singular, so either  $xy_1$  or  $xy_2$  is singular in cl(G). Thus  $y_1y_2 \notin E(cl(G))$ . Let  $G = G_1, \ldots, G_s = cl(G)$  be a locally complete sequence of graphs that yields cl(G). Then  $x \in LD_0(G_i)$  for some *i*. We can deduce that  $x \in LD_0(G_{i+1})$  by the fact that both  $y_1$  and  $y_2$  are not eligible vertices in  $G_i$  and  $y_1y_2 \notin E(G_{i+1})$ . Hence  $x \in LD_0(cl(G))$ , a contradiction. By Lemma 7,  $x \in LD_1(G)$ .  $\Box$ 

The following result is useful for proving our main result.

**Lemma 9.** Let G be a graph satisfying all conditions of Theorem 5. Then cl(G) is a connected claw-free graph such that

- (6) every locally disconnected vertex of degree at least 3 in cl(G) lies on an induced cycle of length at least 4 with at most 3 nonsingular edges;
- (7) every locally disconnected vertex of degree 2 in cl(G) lies on an induced cycle C' with at most 2 non-singular edges such that  $cl(G)[V(C') \cap V_2(cl(G))]$  is a path or a cycle.

In order to prove Lemma 9, we need the following lemmas. A *branch* in *G* is a nontrivial path with end vertices that do not lie in  $V_2(G)$  and with internal vertices of degree 2 (if existing). If a branch has length 1, then it has no internal vertices of degree 2. We use  $\mathscr{B}(G)$  to denote the set of branches in *G*.

**Lemma 10.** Let G be a claw-free graph. If the length of  $L \in \mathcal{B}(G)$  is at least 3 in G, then  $L \in \mathcal{B}(cl(G))$ .

**Proof of Lemma 10.** Let  $G = G_1, \ldots, G_s = cl(G)$  be a locally complete sequence of graphs that yields cl(G). Then *L* is a branch in  $G_i$  for some *i*. Since every vertex of V(L) is not eligible in  $G_i, x_i \in V(G) \setminus V(L)$ , where  $x_i$  is the eligible vertex such that  $G_{i+1} = (G_i)_{x_i}^*$ . Noticing that  $(V(L) \cap V_2(G)) \cap N(x_i) = \emptyset$  and  $|V(L) \cap V_2(G)| \ge 2$ , *L* is a branch in  $G_{i+1}$ . By recursively performing this operation, we can obtain that *L* is a branch of cl(G).  $\Box$ 

**Lemma 11.** Let *G* be a claw-free graph and *C* be an induced cycle with at most *s* non-singular edges in *G* and with at least s - l locally connected vertices in *G*. If  $x \in V(C)$  is locally disconnected in cl(G), then there is an induced cycle *C'* of length at least 4 in cl(G) with  $x \in V(C') \subseteq V(C)$  and with at most *l* non-singular edges in cl(G), where *s* and *l* are nonnegative integers.

**Proof of Lemma 11.** Since  $x \in V(C)$  is locally disconnected in cl(G), there is an induced cycle C' in cl(G) such that  $x \in V(C') \subseteq V(C)$  and  $|V(C')| \ge 4$ . It remains to prove that C' has at most l non-singular edges in cl(G).

Note that every vertex of C' is locally disconnected in cl(G). By Lemma 8,  $V(C') \cap LD_i(cl(G)) \subseteq V(C) \cap LD_i(G)$  for  $i \in \{0, 1, 2\}$ . Hence the number of non-singular edges in C' is no more than s, the number of non-singular edges in C. If C has no locally connected vertex in G, then s = l, hence we are done. Now consider  $s \neq l$ .

Suppose that  $\{u_1, \ldots, u_{s-l}\} \subseteq V(C) \cap LC(G)$ . By Condition (3) in Theorem 6, cl(*G*) is uniquely determined by the graph *G*, i.e., cl(*G*) is independent of the order of eligible vertices during the construction. Note that each  $u_i$  is an eligible vertex in *G* by the hypothesis that *C* is an induced cycle. Let  $G_1 = G_{u_1}^*$  and  $N_G(u_1) \cap V(C) = \{v_1, v_2\}$ . Then there exists an induced cycle

 $C_1$  in  $G_1$  with  $V(C_1) = V(C) \setminus \{u_1\}$  and  $E(C_1) = (E(C) \setminus \{u_1v_1, u_1v_2\}) \cup \{v_1v_2\}$ . Since  $u_1v_1, u_1v_2, v_1v_2$  are non-singular in G,  $C_1$  has at most s - 1 non-singular edges in  $G_1$ . Since  $C_1$  is an induced cycle,  $u_i$  is an eligible vertex in  $G_1$  for  $i \in \{2, \ldots, s - l\}$ . By recursively performing the local completion on  $u_i$  for  $i \in \{1, \ldots, s - l\}$ , we can obtain an induced cycle  $C_{s-l}$  in  $G_{s-l}$ such that  $C_{s-l}$  has at most s - (s - l) = l non-singular edges in  $G_{s-l}$  and  $V(C_{s-l}) = V(C) \setminus \{u_1, \ldots, u_{s-l}\}$ . By Lemma 8,  $(V(C') \cap LD_i(cl(G))) \subseteq (V(C_{s-l}) \cap LD_i(G_{s-l}))$  for  $i \in \{0, 1, 2\}$ . Hence the number of non-singular edges of C' in cl(G) is no more than the number l of non-singular edges of  $C_{s-l}$  in cl(G).  $\Box$ 

Now we provide the proof of Lemma 9.

**Proof of Lemma 9.** First suppose that x is a locally disconnected vertex of degree at least 3 in cl(G). Then either  $x \in LD_1(cl(G))$  or  $x \in LD_2(cl(G))$ . By Lemma 8, either  $x \in LD_1(G)$  or  $x \in LD_2(G)$ . This implies that x is locally disconnected in G and  $d_G(x) \ge 3$ . By Condition (i) of Theorem 5, x lies on an induced cycle of length at least 4 with at most s non-singular edges and with at least s - 3 locally connected vertices. By Lemma 11, x satisfies Condition (6) of Lemma 9.

Next suppose that x is a locally disconnected vertex of degree 2 in cl(G). Then x is a locally disconnected vertex of degree 2 in G. By Condition (ii) of Theorem 5, x lies on an induced cycle C of length at least 4 with at most s non-singular edges and with at least s - 2 locally connected vertices such that  $G[V(C) \cap V_2(G)]$  is a path or a cycle. By Lemma 11, x lies on an induced cycle C' with  $V(C') \subseteq V(C)$  and with at most 2 non-singular edges.

If  $G[V(C) \cap V_2(G)]$  is a cycle, then since *G* is connected, *G* is a cycle. Hence cl(G) is a cycle and we are done. If  $V(C) \cap V_2(G) = \{x\}$ , then since  $x \in V(C') \subseteq V(C)$  and  $x \in V_2(cl(G)) \subseteq V_2(G)$ ,  $V(C') \cap V_2(cl(G)) = \{x\}$  and we are also done. Suppose that  $|V(C) \cap V_2(G)| \ge 2$  and *L* is the branch such that  $(V(C) \cap V_2(G)) \subseteq V(L)$ . By Condition (ii) of Theorem 5,  $L \in \mathcal{B}(G)$  is the unique branch in *C*. By Lemma 10,  $L \in \mathcal{B}(cl(G))$  is the unique branch in *C'*. This implies that  $cl(G)[V(C') \cap V_2(cl(G))]$  is a path (of *L*). Now we complete the proof of Lemma 9.  $\Box$ 

# 3. Proof of Theorem 5

In this section, we present the proof of the main result of this paper. A graph is *Eulerian* if it is connected and has no vertex of odd degree. For a graph *G* with an Eulerian subgraph *H*, we call *H* a spanning Eulerian subgraph of *G* if V(G) = V(H); and a dominating Eulerian subgraph of *G* if G - V(H) is edgeless.

**Theorem 12** (Lai [3]). Let *G* be a 2-connected graph with  $\delta(G) \ge 3$ . If every edge of *G* lies on a cycle of length at most 4, then *G* has a spanning Eulerian subgraph.

Note that the graphs in consideration in Theorem 12 may have multiple edges. A well-known relationship between dominating Eulerian subgraphs in G and Hamiltonian cycles in L(G) was given by Harary and Nash-Williams.

**Theorem 13** (Harary and Nash-Williams [2]). Let G be a graph with at least 3 edges. Then the line graph L(G) is Hamiltonian if and only if G has a dominating Eulerian subgraph.

The following result is immediately from Condition (7) in Lemma 9, which is also necessary for our proof.

**Lemma 14.** Let *G* be a graph satisfying the conditions of Theorem 5. Then every branch  $L \in \mathcal{B}(cl(G))$  of length at least 2 lies on an induced cycle *C* such that *C* has at most 2 non-singular edges and *L* is an unique branch of length at least 2 in *C*.

Before presenting the proofs of main results, we give some additional notation. Let M and M' be the two sets of edges of a graph G. We use  $M \triangle M'$  to denote the symmetric difference of M and M', i.e,  $M \triangle M' = (M \cup M') \setminus (M \cap M')$ . An edge e is called a *pendant* edge if the degree of an end vertex of e is 1; otherwise, it is non-pendant. The graph H for which L(H) = G will be called the *preimage* of G and denote  $H = L^{-1}(G)$ . Note that for subgraph  $G_1 \subseteq G$ ,  $L^{-1}(G_1)$  is possible not unique. However  $L^{-1}(G_1)$  would be unique if  $G_1 \subseteq G$  is an induced subgraph of order at least three. Therefore, for any induced subgraph C of a line graph G, we let  $L^{-1}(C)$  denote the preimage of C.

Given 2-connected block *B* of a simple graph *H* that is not a cycle, let  $U(B) = \{u : d_B(u) = 2 \text{ and } d_H(u) \ge 3\}$  and  $U_1(B) = \{u : u \in U(B) \text{ and } N_B(u) \cap V_2(H) = \emptyset\}$  and  $U_2(B) = U(B) \setminus U_1(B)$ . Now we present the proof of our main result.

**Proof of Theorem 5.** Suppose firstly that *G* is itself a cycle, then it is clearly Hamiltonian and we are done. Now suppose that *G* is not a cycle. By (5) in Theorem 6, it suffices to prove that its closure cl(G) is Hamiltonian. By (4) in Theorem 6, we may assume that cl(G) = L(H) is the line graph of a triangle-free graph *H*.

In order to use Theorem 13, it suffices to find a dominating Eulerian subgraph in *H*. For this, in the following, we use Theorem 12 to get a dominating Eulerian subgraph in each block with some properties and then prove that the union of these Eulerian subgraphs is the desired dominating Eulerian subgraph.

Taking any block B of H that is not a cycle or a pendant edge, we may show the following claims.

**Claim 1.** Every edge  $e = uv \in E(B)$  lies on some cycle C such that either

(8) *C* has exactly 2 vertices of degree greater than 2 in *H* and *C* has exactly one branch of length at least 3 in *H*; or (9) *C* has at most 3 vertices of degree greater than 2 in *H* and *C* has no branch of length at least 3 in *H*.

**Proof of Claim 1.** By Lemma 9, every locally disconnected vertex in cl(G) satisfies Condition (6) or (7) of Lemma 9. Note that if  $e \in E(cl(G))$  is singular, then  $d_H(e) = 2$  because *G* is claw-free. We first prove the following fact.

**Claim 1.0.** Every branch  $L \in \mathcal{B}(H)$  of length at least 3 lies on a cycle C such that C has exactly 2 vertices of degree greater than 2 in H and L is the unique branch of length 3 in C.

**Proof of Claim 1.0.** Because G = L(H), the V(G) is identified with E(H). Thus,  $L' \in \mathcal{B}(cl(G))$  is the line graph of L. Note that  $|L'| = |L| - 1 \ge 2$ . By Lemma 14, there exists an induced cycle C' such that C' has at most 2 non-singular edges and L' is the unique branch of length at least 2 in C'.

By the fact that cl(G) = L(H),  $L^{-1}(C')$  is a cycle in H such that  $L^{-1}(C')$  has at most 2 vertices of degree greater than 2 in H and L is the unique branch of length at least 3 in H. Moreover, since B is not a cycle,  $L^{-1}(C')$  has at least 2 vertices of degree greater than 2 in H. Thus  $L^{-1}(C')$  has exactly 2 vertices of degree greater than 2 in H.  $\Box$ 

Now we start to prove Claim 1. If *e* lies on a branch  $L \in \mathcal{B}(H)$  of length at least 3, then *e* lies on a cycle satisfying (8) by Claim 1.0. Now suppose that *e* lies on a branch  $L \in \mathcal{B}(H)$  of length 1 or 2. Let  $v_e \in V(cl(G))$  be the vertex corresponding to the edge *e* in *H*. Then  $d_{cl(G)}(v_e) \ge 3$ . Since *H* is triangle free, we have  $N_H(u) \cap N_H(v) = \emptyset$ . By the fact that *B* is 2-connected,  $N_H(u) \neq \emptyset$  and  $N_H(v) \neq \emptyset$ . Furthermore, since cl(G) is claw-free,  $cl(G)[N_{cl(G)}(v_e)]$  is composed of two vertex-disjoint cliques, i.e.,  $v_e$  is locally disconnected in cl(G).

By the hypotheses of Lemma 9,  $v_e$  lies on an induced cycle  $C_e$  of length at least four with at most 3 non-singular edges. Hence, by the fact that  $cl(G) = L(H), L^{-1}(C_e)$  is a cycle in H such that  $e \in E(L^{-1}(C_e))$  and  $|V(L^{-1}(C_e)) \cap V_{\geq 3}(H)| \leq 3$ . Since B is not a cycle,  $|V(L^{-1}(C_e)) \cap V_{\geq 3}(H)| \geq 2$ . Note that  $L^{-1}(C_e)$  has a branch of length 1 or 2. Therefore, if  $|V(L^{-1}(C_e)) \cap V_{\geq 3}(H)| = 2$ , then  $L^{-1}(C_e)$  has at most one branch of length at least 3 which implies that  $L^{-1}(C_e)$  satisfies (8); if  $|V(L^{-1}(C_e)) \cap V_{\geq 3}(H)| = 3$  and  $L^{-1}(C_e)$  has t branches of length at least 3, then  $t \leq 2$ . Suppose that  $L_1, L_2$  are the two possible branches of length at least 3 in  $L^{-1}(C_e)$  ( $L_1 = L_2$  if t = 1). By Claim 1.0,  $L_i$  lies on a cycle  $C_i$  satisfying (8) for i = 1, 2. Thus, either  $C = H[E(L^{-1}(C_e) \triangle E(C_1)) \triangle E(C_2)]$  (if t = 2), or  $C = H[E(L^{-1}(C_e)) \triangle E(C_1)]$  (if t = 1), or  $C = C_e$  (if t = 0) is a cycle such that  $e \in E(C)$  and such that C satisfies (9), which completes the proof of Claim 1.

**Claim 2.** Every vertex u of U(B) lies on a cycle C such that C satisfies (9) of Claim 1 and  $|V(C) \cap V_{\geq 3}(B)| = 2$  and  $V(C) \cap U(B) = \{u\}$ .

**Proof of Claim 2.** Let *e* be an edge incident with *u* in *B*. By Claim 1, *e* lies on a cycle *C* satisfying Condition (8) or (9) of Claim 1. This implies that  $|V(C) \cap V_{\geq 3}(H)| \leq 3$ . Since *B* is not a cycle,  $|V(C) \cap V_{\geq 3}(B)| \geq 2$ . Note that  $u \in V(C) \cap (V_{\geq 3}(H) \setminus V_{\geq 3}(B))$ ,  $|V(C) \cap V_{\geq 3}(H)| \geq 3$ . Thus  $|V(C) \cap V_{\geq 3}(H)| = 3$  and  $|V(C) \cap V_{\geq 3}(B)| = 2$ . This implies that *C* satisfies Condition (9) of Claim 1 and  $V(C) \cap U(B) = \{u\}$ .  $\Box$ 

**Claim 3.** Every vertex *u* of  $U_1(B)$  lies on a cycle *C* of length 4 such that  $|V(C) \cap V_2(B)| = 2$ ,  $|V(C) \cap V_{\geq 3}(H)| = 3$  and  $|V(C) \cap V_2(H)| = 1$ .

**Proof of Claim 3.** By Claim 2, we may take a cycle *C* such that *C* satisfies Condition (9) of Claim 1. Note that  $N_B[u] \subseteq V(C)$  and  $N_B[u] \subseteq V_{\geq 3}(H)$ ,  $|V(C) \cap V_{\geq 3}(H)| = 3$ . Since *H* is triangle-free and there is no branch of length at least 3 in *C*, |V(C)| = 4 and  $|V(C) \cap V_2(H)| = 1$ . Since  $u \in V_2(B) \setminus V_2(H)$ ,  $|V(C) \cap V_2(B)| = 2$ .  $\Box$ 

Using the above claims, we obtain three kinds of blocks  $B_1$ ,  $B_2$ ,  $B_3$  from B by the following three steps. Note that possibly  $B_i = B_{i-1}$  for some  $\leq 2$ , where  $B_0 = B$ .

(I) Let  $B_1$  be obtained from B by replacing each branch of length 2 with an edge. Then we have the following fact.

**Fact 1.** If  $B_1$  has a spanning Eulerian subgraph, then B has a dominating Eulerian subgraph F such that  $(V_3(B) \cup (V_{\geq 3}(H) \cap V(B))) \subseteq V(F)$ , i.e.,  $(V_3(B) \cup U(B)) \subseteq V(F)$ .

Clearly  $V_1(H) \cap V(F) = \emptyset$ . Furthermore, there may be some vertices in  $V_2(B) \cap V_2(H)$  that are not in V(F), i.e., possibly  $(V_2(B) \cap V_2(H)) \setminus V(F) \neq \emptyset$ . However, Fact 1 shows that  $(V_{\geq 3}(H) \cap V_2(B)) \setminus V(F) = \emptyset$ .

*Proof of* Fact 1. Suppose  $B_1$  has a spanning Eulerian subgraph  $F_1$ . By the definition of  $B_1$  in Step (I), B has a dominating Eulerian subgraph F such that  $V_{\geq 3}(B) \subseteq V(F_1) \subseteq V(F)$ . Note that  $(V_{\geq 3}(H) \cap V(B)) \setminus V_{\geq 3}(B) = U(B)$ . Since any  $u \in U_2(B)$  lies on a branch of length at least 3 in B,  $U_2(B) \subseteq V(F)$  by the definition of  $B_1$  in Step (I). Now suppose  $u \in U_1(B)$ , then u lies on a branch of length 2 in B. Furthermore, by Claim 3, u lies on a 4-cycle C with two branches of length 2 in B with  $|V(C) \cap V_2(B)| = 2$ ,  $|V(C) \cap V_{\geq 3}(H)| = 3$  and  $|V(C) \cap V_2(H)| = 1$  and there would exist two parallel edges in  $B_1$  by Step (I). Note that one of the parallel edges is obtained by replacing branch containing u and the other is



**Fig. 2.** The graphs  $G_5$  and  $G_6$ .

obtained by replacing the branch of length 2 containing the vertex of  $V(C) \cap V_2(H)$ . Note also that  $F_1$  may be chosen such that it contains at least one of the parallel edges. Hence we may assume that  $F_1$  is a spanning Eulerian subgraph of  $B_1$  such that the edge obtained by replacing the branch containing u belong to  $F_1$ . This implies that  $U_1(B) \subseteq V(F)$ . Thus,  $(V_{\geq 3}(H) \cap V(B)) \subseteq V(F)$ . By the definition of  $B_1, V_3(B) \subseteq V(F)$ . This completes the proof of Fact 1.  $\Box$ 

(II) Let  $B_2$  be obtained from  $B_1$  by replacing each branch *L* of length at least 3 with a branch of length 3. Then

**Fact 2.** If  $B_2$  has a spanning Eulerian subgraph, then  $B_1$  has a spanning Eulerian subgraph.

(III) Let  $v, v' \in V(L) \cap V_2(B_2)$  and let  $B_3$  be a graph obtained from  $B_2$  and two additional vertices  $w_v, w_{v'}$  by adding all the edges  $\{w_v v, w_v v', w_v v', w_{v'} v', w_v w_{v'}\}$ . Then  $vv' \in E(B_2)$  and

**Fact 3.** If  $B_3$  has a spanning Eulerian subgraph, then  $B_2$  has a spanning Eulerian subgraph.

By the definition of *B*<sub>3</sub>, and by Claims 1 and 2, we may obtain the following fact immediately.

**Fact 4.** Every edge of  $B_3$  lies in a cycle of length at most 4 and  $\delta(B_3) \ge 3$ .

Note that the vertex of degree greater than 2 is the dominating Eulerian subgraph of a block that is a pendant edge and H has no block of order 2 that is not a pedant edge by Lemma 9. Hence, since  $B_3$  is 2-connected,  $B_3$  has an spanning Eulerian subgraph by Fact 4 and by Theorem 12. Combining Facts 1–3, one may obtain the following claim.

**Claim 4.** Every block *B* of *H* that is not a cycle has a dominating Eulerian subgraph containing all vertices of  $V_3(B) \cup (V_{\geq 3}(H) \cap V(B))$ .

Note that every block of *H* that is a cycle has a spanning Eulerian subgraph which is itself the cycle. This together with Claim 4 implies that the union of those dominating Eulerian subgraphs of these blocks is a dominating Eulerian subgraphs of *H*. (Note that every cut vertex of *H* has degree at least three in *H* and then, by Claim 4, it lies on a dominating Eulerian subgraph of the nontrivial blocks containing it. Thus, the union of these dominating Eulerian subgraphs of these blocks is connected by its definition.) By Theorem 13, cl(G) = L(H) is Hamiltonian. The proof of Theorem 5 is complete.

# 4. Proof of Theorem 3 – Theorem 5 implies Theorem 3

Before proving Theorem 3, we start with the following lemmas.

**Lemma 15.** Let *G* be a connected claw-free graph satisfying Condition (1) of Theorem 2. If *G* has an induced subgraph  $H \cong G_5$  such that every edge of *H* is non-singular in *G* and such that *H* has a locally disconnected vertex *x* of *G* with  $d_H(x) = 4$  that does not satisfy Condition (i) of Theorem 5, then  $G \cong G_6$  (see Fig. 2).

**Proof of Lemma 15.** Let *H* be an induced subgraph of *G* isomorphic to the graph  $G_5$  in Fig. 2. If *H* has a locally connected vertex in *G*, then every locally disconnected vertex *x* of *H* with  $d_H(x) = 4$  satisfies Condition (i) of Theorem 5, contradicting the hypothesis. Now suppose that

(\*) every vertex of *H* is locally disconnected in *G*.

All the subscripts are taken module by 4 in the whole proof.

**Claim 5.** For any  $x \in V(G - H)$ , if  $\{xv_i, xv_{i+2}\} \subset E(G)$  for some  $i \in \{1, 2\}$ , then  $xu_j \notin E(G)$  for any  $j \in \{1, 2, 3, 4\}$ .

**Proof of Claim 5.** By contradiction. Without loss of generality, we may assume that there exists a vertex  $x_0 \in V(G-H)$  such that  $\{x_0v_1, x_0v_3\} \subset E(G)$  and  $x_0u_1 \in E(G)$ . Hence  $\{v_1, v_3, u_1\} \subseteq N(x_0)$  and  $G[\{v_1, v_3, u_1\}]$  is connected but not complete by the hypothesis that H is an induced subgraph of G. This implies that  $x_0$  is an eligible vertex of G.

Noticing that  $u_1$  is locally disconnected in G and  $\{u_2, v_1, x_0\}$  belongs to a connected component of  $G[N(u_1)]$ , we have that  $x_0u_2 \in E(G)$  and  $x_0u_4 \notin E(G)$ . Since  $u_2$  is locally disconnected,  $x_0u_3 \notin E(G)$ . Hence we can obtain two induced cycles  $C_1 = u_1u_4v_3x_0u_1$  and  $C_2 = u_2u_3v_3x_0u_2$  such that they have exactly four non-singular edges and have a locally connected vertex  $x_0$  of G. This implies that every locally disconnected vertex of degree 4 in H satisfies (i) of Theorem 5, contradicting the hypothesis.  $\Box$ 

**Claim 6.** Let  $C_1$  and  $C_2$  be any two triangles containing  $v_1v_3$  and  $v_2v_4$ , respectively. Then  $|V(C_1) \cap V(C_2)| = 1$ .

**Proof of Claim 6.** Suppose, otherwise, that  $C_1 = w_1v_1v_3w_1$  and  $C_2 = w_2v_2v_4w_2$  and  $w_1 \neq w_2$ . Since  $H \cong G_5$  is an induced subgraph,  $\{w_1, w_2\} \subseteq V(G - H)$ . Consider  $S_j = (V(H) \setminus \{v_{j+1}\}) \cup \{w_j\}$  for  $j \in \{1, 2\}$ . By Claim 5,  $w_ju_i \notin E(G)$  for any  $i \in \{1, 2, 3, 4\}$ . Suppose  $w_jv_{j+3} \notin E(G)$ , then  $G[S_j] \cong G_2$ . By the hypothesis of Lemma 15, there is at least a locally connected vertex in  $V(H) \setminus \{v_{j+1}\}$ . This contradicts (\*). Thus  $w_jv_{j+3} \in E(G)$ . By symmetry, we can obtain  $w_jv_{j+1} \in E(G)$ .

Since  $\{u_1, u_2, w_1, w_2\} \subset N(v_1), w_1w_2 \in E(G)$ : otherwise  $G[\{v_1, w_1, w_2, u_1\}]$  would be a claw by Claim 5, a contradiction. Hence  $\{w_2, v_1, v_2, v_3, v_4\} \subseteq N(w_1)$  and  $G[\{w_2, v_1, v_2, v_3, v_4\}]$  is connected but not complete. This implies that  $w_1$  is an eligible vertex. Hence we can obtain that  $C_i = u_i v_i w_1 v_{i-1} u_i$  is an induced cycle with a locally connected vertex  $w_1$  that implies  $u_i$  satisfying Condition (i) of Theorem 5 for any  $i \in \{1, 2, 3, 4\}$ , contradicting the hypothesis.  $\Box$ 

By Claim 6, we may assume that w is the common neighbor of  $\{v_1, v_2, v_3, v_4\}$ . Noticing that H is an induced subgraph in G and  $wu_i \notin E(G)$  by Claim 5, we can obtain that  $G[V(H) \cup \{w\}] \cong G_6$ . Now suppose that  $x \in V(G) \setminus (V(H) \cup \{w\})$ . If  $xw \in E(G)$ , then since  $\{v_1, v_2, v_3, v_4, x\} \subseteq N(w)$  and by the fact that G is claw-free, then  $xv_i, xv_{i+2} \in E(G)$  for some  $i \in \{1, 2\}$ , say i = 1. Hence we can obtain two triangles  $xv_1v_3x$  and  $wv_2v_4w$ . By Claim 6, x = w. However, this contradicts that  $x \in V(G) \setminus (V(H) \cup \{w\})$ . Hence  $xw \notin E(G)$ .

If  $xu_i \in E(G)$  for some  $i \in \{1, 2, 3, 4\}$ , then since  $\{u_{i+1}, u_{i-1}, v_{i-1}, v_i, x\} \subset N(u_i)$  and by the fact that *G* is claw-free,  $\{xv_i, xu_{i+1}\} \subset E(G)$  or  $\{xv_{i-1}, xu_{i-1}\} \subset E(G)$ , say  $\{xv_i, xu_{i+1}\} \subset E(G)$ . Consider  $S' = (V(H) \setminus \{v_i\}) \cup \{x\}$ . By Claim 5,  $xv_{i+2} \notin E(G)$ . Since  $u_i, u_{i+1}$  are locally disconnected,  $\{xu_{i+2}, xu_{i+3}, xv_{i-1}, xv_{i+1}\} \cap E(G) = \emptyset$ . Thus  $G[S'] \cong G_2$ . By the hypothesis of Lemma 15, there is a locally connected vertex in  $V(H) \setminus \{v_i\}$ , contradicting (\*). Thus  $xu_i \notin E(G)$ . Similarly, we can obtain  $xv_i \notin E(G)$  for any  $i \in \{1, 2, 3, 4\}$ . Since *G* is connected,  $V(G) = V(H) \cup \{w\}$ . Thus  $G \cong G_6$ . We have now completed the proof of Lemma 15.  $\Box$ 

**Lemma 16.** Let *G* be a claw-free graph satisfying Condition (1) of Theorem 2. If a locally disconnected  $u_1$  lies on an induced cycle of length 4, then either  $u_1$  satisfies Condition (i) or (ii) of Theorem 5; or  $G \cong G_6$ .

**Proof of Lemma 16.** Suppose that *C* is an induced cycle of length 4 with  $u_1 \in V(C)$ . If either  $V(C) \cap LC(G) \neq \emptyset$  or *C* has a singular edge, then  $u_1$  satisfies Condition (i) or (ii) of Theorem 5 (Note that if  $d(u_1) = 2$ , then the two edges incident with  $u_1$  are singular edges and clearly  $u_1$  satisfies Condition (ii) of Theorem 5.) and we are done.

Now consider the case when  $V(C) \cap LC(G) = \emptyset$  and *C* has no singular edge. Let  $C = u_1u_2u_3u_4u_1$  and  $T_i = v_iu_iu_{i+1}v_i$  be a triangle containing  $u_iu_{i+1}$ , and let  $S = \bigcup_{i=1}^{i=4} V(T_i)$ , where  $i \in \{1, 2, 3, 4\}$  and all the subscripts are taken module by 4 in the proof. Note that  $v_iv_{i+1} \notin E(G)$  and  $v_i \neq v_{i+1}$  since  $u_{i+1}$  is locally disconnected. Similarly,  $v_1, v_2, v_3, v_4$  are four distinct vertices, since, otherwise, say,  $v_1 = v_3$ , then  $u_1$  would be locally connected, a contraction. Since *C* is an induced cycle in *G* and every vertex of *C* is locally disconnected in *G*,  $\{u_iv_{i+1}, u_iv_{i+2}, v_iv_{i+1} : i \in \{1, 2, 3, 4\} \cap E(G) = \emptyset$ .

Suppose first that  $\{v_1v_3, v_2v_4\} \cap E(G) = \emptyset$ , then  $G[S] \cong G_1$ . By the hypothesis of Theorem 2, some  $u_i \neq u_1$  is locally connected in *G*. Hence  $u_1$  satisfies Condition (i) of Theorem 5.

Suppose next that exactly one of  $\{v_1v_3, v_2v_4\}$ , say  $v_1v_3$ , belongs to E(G), then  $G[S] \cong G_2$ . By the hypothesis of Theorem 2, at least one of  $\{u_1, u_2, u_3, u_4, v_1, v_3\}$  is locally connected in *G*. Hence  $u_1$  satisfies Condition (i) of Theorem 5.

Suppose finally that  $\{v_1v_3, v_2v_4\} \subset E(G)$ . Then  $G[S] \cong G_5$ . If exactly one of  $\{v_1v_3, v_2v_4\}$ , say  $v_2v_4$ , is non-singular in G, then we can obtain an induced cycle  $C' = u_1v_1v_3u_4u_1$  such that  $u_1$  satisfies Condition (i) of Theorem 5. Now suppose that both  $v_1v_3$  and  $v_2v_4 \in E(G)$  are non-singular in G. If  $u_1$  satisfies Condition (i) of Theorem 5, then we are done. Otherwise by Lemma 15,  $G \cong G_6$  which is Hamiltonian. This completes the proof of Lemma 16.  $\Box$ 

#### Now we prove Theorem 3.

**Proof of Theorem 3.** We shall prove that the locally disconnected vertices in *G* satisfy Condition (i) or (ii) of Theorem 5 with only one exceptional case when  $G \cong G_6$ . Suppose that *u* is a locally disconnected vertex of *G*. Since *G* is  $N^2$ -locally connected, *u* lies on an induced cycle *C* of length 4 or 5. If |V(C)| = 4, then by Lemma 16, either *u* satisfies Condition (i) or (ii) of Theorem 5 or  $G \cong G_6$  and we are done.

Now suppose |V(C)| = 5. If either  $|V(C) \cap LC(G)| \ge 2$  or *C* has at most 2 non-singular edges, then *u* satisfies Condition (i) or (ii) of Theorem 5 (Note that if *C* has a vertex of degree two, then *C* has at most three non-singular edges.) and we are done. It remains to consider the case when  $|V(C) \cap LC(G)| \le 1$  and *C* has at least 3 non-singular edges. Let  $C = u_1u_2u_3u_4u_5u_1$ , where  $u = u_i$  for some *i*. All the subscripts are taken by 5 in the whole proof.

**Case 1.** *C* has a path *P* of length 3 such that every vertex of *P* is locally disconnected and every edge of *P* is non-singular.

Without loss of generality, we may assume that  $P = u_2 u_3 u_4 u_5$ . Let  $T_i = v_i u_{i+1} u_{i+2} v_i$  be a triangle for i = 1, 2, 3, and let  $S = V(C) \bigcup \bigcup_{i=1}^{i=3} V(T_i)$ . Since each vertex of P is locally disconnected,  $v_1, v_2, v_3$  are distinct vertices. Consider the induced subgraph G[S]. Since each  $u_i \neq u_1$  is locally disconnected in G,  $\{v_1v_2, v_2v_3, v_iu_i, v_iu_{i+3} : i \in \{1, 2, 3\}\} \cap E(G) = \emptyset$ . Thus  $v_iu_{i-1} \notin E(G)$ : Otherwise,  $G[\{u_{i-1}, v_i, u_i, u_{i-2}\}]$  is a claw since C is an induced cycle, a contradiction. If  $v_1v_3 \notin E(G)$ , then this implies that  $G[S] \cong G_3$ ; this is impossible by Condition (2) of Theorem 3. Thus  $v_1v_3 \in E(G)$ , implying that  $G[S] \cong G_4$  and hence at least one of  $\{v_1, v_3\}$  is locally connected by Condition (2) of Theorem 3.

First suppose that either  $u_1$  is locally connected or at least one of  $\{u_1u_5, u_1u_2\}$  is singular, then u lies on either an induced cycle  $C' = u_1u_2v_1v_3u_5u_1$  or an induced cycle  $C'' = v_1u_3u_4v_3v_1$ , hence u satisfies Condition (i) or (ii) of Theorem 5 and we are done.

Next suppose that  $u_1$  is locally disconnected and both  $u_1u_5$  and  $u_1u_2$  are non-singular, then let  $T_i = v_iu_{i+1}u_{i+2}v_i$  be a triangle for  $i \in \{4, 5\}$ , and let  $S' = V(C) \cup \bigcup_{i=3}^{i=5} V(T_i)$ . Then  $v_4 \neq v_5$ . Now consider the induced subgraph G[S']. Similar to the discussion on G[S] of the fact that  $v_1v_3 \in E(G)$ , we can obtain that  $v_3v_5 \in E(G)$ . But this implies that  $G[\{v_3, u_4, v_1, v_5\}]$  is a claw, a contradiction.

**Case 2.** *C* has no a path *P* of length 3 such that every vertex of *P* is locally disconnected and every edge of *P* is non-singular.

If *C* has no locally connected vertex, then it suffices to consider the case when *C* has exactly 3 non-singular edges of *G* which are not consecutive. This implies that the degree of every vertex of *C* is at least 3 in *G*. Hence *u* satisfies Condition (i) of Theorem 5.

If *C* has a locally connected vertex, then it suffices to consider the case when *C* has at least one singular edge in *G*. If  $d(u) \ge 3$ , then *C* has at most 4 non-singular edges and a locally connected vertex. If d(u) = 2, then *C* has at most 3 non-singular edges and a locally connected vertex. This implies that *u* satisfies Condition (i) or (ii) of Theorem 5, in either case above. This completes the proof of Theorem 3.  $\Box$ 

#### 5. Concluded remarks

#### 5.1. Sharpness

In this subsection, we discuss the sharpness and show that the conditions of Theorems 5 and 3 are all best possible in some sense.

- Let  $k \ge 2$  be an integer. Let  $G^k$  be the graph obtained from  $K_{2,3}$  by attaching k pendant edges to each vertex of  $K_{2,3}$ , and then subdividing each original edge once. It is straightforward to check that the graph  $H^k = L(G^k)$  is 2-connected claw-free graph with  $\delta(H^k) \ge k + 1$  such that every locally disconnected vertex of G lies on an induced cycle of length 8 with 4 non-singular edges and without a locally connected vertex. However, it is not Hamiltonian: Otherwise, by Theorem 13,  $G^k$  has a dominating Eulerian subgraph, a contradiction. This shows that "s 3 locally connected vertices" in Condition (i) of Theorem 5 cannot be replaced by "s 4 locally-connected vertices" even under an additional condition that a graph has any given large minimum degree.
- We demonstrate that the condition " $G[V(C) \cap V_2(G)]$  is a path or a cycle" in Theorem 5(ii) is necessary. The graph G' in Fig. 3 is a graph satisfying Condition (i) but not (ii) since the two locally disconnected vertices of degree two do not satisfy Condition (ii) of Theorem 5 (although they lie on an induced cycle with only two nonsingular edges). It is straightforward to check that G' is not Hamiltonian. This shows that the condition " $G[V(C) \cap V_2(G)]$  is a path or a cycle" in Condition (ii) of Theorem 5 is necessary. One can obtain many such graphs with any large order by joining a clique of any large order to any nontrivial maximal clique of G'.
- Theorem 3 is best possible in the sense that Condition (2) cannot be weakened. To see this, consider G' depicted in Fig. 3. The graph G' satisfies all conditions of Theorem 3 except that every induced subgraph H of G' isomorphic to  $G_3$  does not have a locally connected vertex in  $V_3(H) \cup V_4(H)$ . Note that G' is not Hamiltonian. This shows that "at least one locally connected vertex" in Condition (2) of Theorem 3 is necessary.

#### 5.2. A flaw in the original proof of Corollary 4

Looking at the original proof of Corollary 4, we find that the author used the assumption that if a claw-free graph satisfies the condition that it is connected  $N^2$ -locally connected  $\{G_1, G_2, G_3, G_4\}$ -free (where  $G_i$  is the graph in Fig. 1), then so does its closure. However, this is not generally true: For example, the claw-free graph G'' in Fig. 3 is connected  $N^2$ -locally connected  $\{G_1, G_2, G_3, G_4\}$ -free, however, cl(G'') has an induced subgraph isomorphic to  $G_3$ . This implies that cl(G'') does not satisfy the conditions of Corollary 4 any more. One can obtain many such graphs of any large order by joining a clique of any large order to any nontrivial maximal clique of G''. Our proof has conquered the flaw.

# 5.3. Theorem 5 is stronger than Theorem 3 and Corollary 4

We show that some graphs satisfy Theorem 5 but not Theorem 3. As we showed in Section 5.2, the closure of the graph G'' depicted in Fig. 3 does not satisfy the conditions of Corollary 4. In fact, it does not satisfy Theorem 3. However, it satisfies the conditions of Theorem 5. This shows that Theorem 5 is stronger than Theorem 3.



**Fig. 3.** The graph G', G'': *u* is a locally connected vertex of G''.

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