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# Improving results on the pseudorandomness of sequences generated via the additive order of a finite field



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#### 1. Introduction

## A B S T R A C T

We improve several results in the area of pseudorandom sequences. First, we obtain an improved bound on the general lattice test for digital explicit inversive and digital explicit nonlinear pseudorandom number generators. Second, we improve the bound on the correlation measure of binary sequences generated by the quadratic character of finite fields. Finally, we improve the bound on the correlation measure of digital explicit inversive pseudorandom numbers, and the bound on their linear complexity profile.

Although we follow essentially the earlier proofs, we improved a crucial step, namely a better estimate on the number of nonempty intersections of 'boxes' of a finite field is given. © 2015 Elsevier B.V. All rights reserved.

Let  $q = p^r$  be a prime power and  $\mathbb{F}_q$  be the finite field with q elements. We identify the finite field  $\mathbb{F}_p$  with the set of integers  $\{0, 1, \ldots, p-1\}$ . Let  $\beta_1, \ldots, \beta_r \in \mathbb{F}_q$  be a basis of  $\mathbb{F}_q$  over  $\mathbb{F}_p$ . We define the additive order of  $\mathbb{F}_q$  in the following way: for  $n \in \{0, 1, \ldots, q-1\}$  let

$$\xi_n = n_1\beta_1 + n_2\beta_2 + \dots + n_r\beta_r$$

if

 $n = n_1 + n_2 p + \dots + n_r p^{r-1}, \quad 0 \le n_1, n_2, \dots, n_r < p.$ 

We define  $\xi_{n+q} = \xi_n$  for  $n \in \{0, 1, ..., q-1\}$ . Let  $W \subseteq \mathbb{F}_q$  be defined as follows:

 $\mathcal{W} = \{w_2\beta_2 + \dots + w_r\beta_r : w_2, \dots, w_r \in \{0, 1\}\}.$ 

For an element  $\omega \in W$  and  $a \in \{0, 1, ..., q - 1\}$  we define

 $S_{a,\omega} = \{\xi_n : 0 \le n < q, \xi_{n+a} = \xi_n + \xi_a + \omega\}.$ 

Let  $d_1, \ldots, d_k$  be integers with  $0 \le d_1 < \cdots < d_k < q$ . We look for an upper bound on the number of nonempty elements in the following set of  $\mathbb{F}_q$ 

 $\{S_{d_1,\omega_1}\cap\cdots\cap S_{d_k,\omega_k}:\omega_1,\ldots,\omega_k\in \mathcal{W}\}.$ 

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Before stating the main result, we give the definition of a box and a remark on the set  $S_{a,\omega}$ . Let  $N_{i,1}$  and  $N_{i,2}$  be integers such that  $0 \le N_{i,1} < N_{i,2} < p$ . We call a set of the form

 $\{\ell_1 \beta_1 + \dots + \ell_r \beta_r : N_{i,1} \le \ell_i < N_{i,2}, i = 1, 2, \dots, r\}$ 

as a box. Let  $w_1 = 0$ ,  $\omega = w_1\beta_1 + \dots + w_r\beta_r \in W$  and  $a = a_1 + a_2p + \dots + a_rp^{r-1}$  for some  $a_1, a_2, \dots, a_r \in \{0, 1, \dots, p-1\}$ . Then  $S_{a,\omega}$  has the form

$$S_{a,\omega} = \left\{ \ell_1 \beta_1 + \ell_2 \beta_2 + \dots + \ell_r \beta_r : \max\{0, pw_{i+1} - a_i - w_i\} \le \ell_i < \min\{p, pw_{i+1} - a_i - w_i + p\}, \\ i = 1, 2, \dots, r - 1, \ 0 \le \ell_r < p \right\}.$$

Hence,  $S_{a,\omega}$  is a box. We now state our main theorem.

**Theorem 1.** Let  $d_1, \ldots, d_k$  be integers with  $0 \le d_1 < \cdots < d_k < q$  and let  $\omega_1, \ldots, \omega_k \in W$ . Then the set

 $S_{d_1,\omega_1} \cap \dots \cap S_{d_k,\omega_k} = \{\xi_n : 0 \le n < q, \xi_{n+d_i} = \xi_n + \xi_{d_i} + \omega_i, i = 1, 2, \dots, k\}$ (2)

is a box or an empty set. If  $\omega_1, \ldots, \omega_k$  run over W, then there are at most  $(k+1)^{r-1}$ -many nonempty sets among them.

Using Theorem 1 one can obtain a similar result in the incomplete case.

**Corollary 1.** Let  $M \in \{0, 1, ..., q-1\}$  and  $\mathcal{E} = \{\xi_0, \xi_1, ..., \xi_{M-1}\} \subseteq \mathbb{F}_q$ . Let  $d_1, ..., d_k$  be integers with  $0 \le d_1 < \cdots < d_k$  $d_k < q$  and let  $\omega_1, \ldots, \omega_k \in W$ . Then, the set

$$S_{d_1,\omega_1} \cap \dots \cap S_{d_k,\omega_k} \cap \mathcal{E} = \{\xi_n : 0 \le n < M, \xi_{n+d_i} = \xi_n + \xi_{d_i} + \omega_i, i = 1, 2, \dots, k\}$$
(3)

can be split into a union of boxes, such that if  $\omega_1, \ldots, \omega_k$  run over W, then the number of boxes is  $O((k+1)^{r-1})$ .

Before presenting the proofs of theorem and corollary, we give some of their applications in Section 2. In particular, we improve the results given in [4,10,11], and [3] in Sections 2.1–2.3 respectively. Next, in Section 3 we give the proofs of Theorem 1 and Corollary 1.

#### 2. Applications

In this section we apply Theorem 1 and Corollary 1 to obtain better results on pseudorandomness of certain sequences.

#### 2.1. On the lattice structure of digital explicit inversive and nonlinear generators

Let

$$\overline{\gamma} = \begin{cases} \gamma^{-1} & \text{if } \gamma \in \mathbb{F}_q^*, \\ 0 & \text{if } \gamma = 0. \end{cases}$$

For given  $\alpha \in \mathbb{F}_q^*$  and  $\beta \in \mathbb{F}_q$ , a sequence  $\gamma_0, \gamma_1, \ldots$  generated by

$$\gamma_n = \alpha \xi_n + \beta, \quad n = 0, 1, \dots$$

is called digital explicit inversive pseudorandom number generator (or Niederreiter–Winterhof generator), see [7]. The inversive pseudorandom number generator is a special case of digital explicit nonlinear pseudorandom number generators  $(n_n)$ defined by

$$\eta_n = f(\xi_n) \tag{5}$$

for some polynomial  $f(X) \in \mathbb{F}_q[X]$ , see [8]. Note that for the inversive generator we have  $f(X) = (\alpha X + \beta)^{q-2}$ . In this section we study the general lattice test (first introduced by Niederreiter and Winterhof [9]) for the digital explicit

inversive and nonlinear generators. Let  $(\eta_n)$  be a *T*-periodic sequence over  $\mathbb{F}_q$ . For given integers  $s \ge 1, 0 < d_1 < d_2 < \cdots < d_n <$  $d_{s-1} < T$ , and  $N \ge 2$ , we say that  $(\eta_n)$  passes the *s*-dimensional *N*-lattice test with lags  $d_1, d_2, \ldots, d_{s-1}$  if the vectors

$$\{\eta_n - \eta_0 : 1 \le n < N\}$$

span  $\mathbb{F}_q^s$ , where

$$\eta_n = (\eta_n, \eta_{n+d_1}, \dots, \eta_{n+d_{s-1}}), \quad 0 \le n < N.$$

The greatest dimension s such that  $(\eta_n)$  satisfies the s-dimensional N lattice test for all lags  $d_1, \ldots, d_{s-1}$  is denoted by  $S(\eta_n, N).$ 

(4)

Pirsic and Winterhof [10] studied the lattice structure of inversive and nonlinear pseudorandom number generators. They used the upper bound  $2^{k(r-1)}$  on the number of possible boxes (2), and obtain that

$$S(\gamma_n, N) \ge \frac{\log(N) - \log\log(N) - 1}{r - 1} - 1$$
 (6)

and

$$S(\eta_n, N) \ge \frac{\log(N/D)}{r-1} - 1$$

for  $2 \le N \le q$  and  $r \ge 2$ , where  $(\gamma_n)$  is defined by (4), and  $(\eta_n)$  is defined by (5) with a nonidentical zero function  $f(X) \in \mathbb{F}_q[X]$  of degree *D*.

Later Gómez-Pérez and Gómez [4] improved the bound (6)

$$S(\gamma_n, N) \ge \frac{1}{6} \left(\frac{N}{r^{r-1}}\right)^{1/r} \tag{7}$$

for  $2 \le N \le q$  and  $r \ge 2$ .

Using Theorem 1 in the proof of [10] gives the following improved bounds:

$$S(\gamma_n, N) \ge \left(\frac{N}{2}\right)^{1/r} - 1$$

and

$$S(\eta_n, N) \ge \left(\frac{N}{D}\right)^{1/(r-1)} - 1$$

for  $2 \le N \le q$  and  $r \ge 2$ . We remark that in these bounds the constant terms do not depend on *r* anymore.

2.2. On the correlation measure of binary sequences defined by the quadratic character

Let

$$E_N = \{e_0, \ldots, e_{N-1}\} \in \{-1, +1\}^N$$

be a binary sequence of length N. The correlation measure of order k of  $E_N$  is defined as

$$C_k(E_N) = \max_{M,D} \left| \sum_{n=0}^{M-1} e_{n+d_1} e_{n+d_2} \dots e_{n+d_k} \right|,$$

where the maximum is taken over all  $D = (d_1, \ldots, d_k)$  and M such that  $0 \le d_1 < d_2 < \cdots < d_k \le N - M$ . This measure was first introduced by Mauduit and Sárközy [5]. For a "good" pseudorandom sequence  $E_N$ ,  $C_k(E_N)$  (for "small" k) is small and is ideally greater than  $N^{1/2}$  only by at most a power of log N, see [1].

The linear complexity profile is an important cryptographic characteristic of pseudorandom sequences. A low linear complexity profile has turned out to be undesirable for cryptographic applications.

We recall that the *linear complexity profile*  $L(R_T, N)$  of the sequence  $R_T = (r_0, r_1, \ldots, r_{T-1})$  is a non-decreasing sequence where the *N*th term is defined as the shortest length *L* of a linear recurrence relation over  $\mathbb{F}_2$ 

 $r_{n+L} = c_{L-1}r_{n+L-1} + \cdots + c_0r_n, \quad 0 \le n \le N - L - 1,$ 

which is satisfied by this sequence.

Let  $\chi$  be the quadratic character of  $\mathbb{F}_q$ . Sárközy and Winterhof [11] defined a binary sequence  $L_q = (l_0, l_1, \dots, l_{q-1}) \in \{-1, 1\}^q$  by

$$l_n = \begin{cases} \chi(f(\xi_n)) & \text{if } f(\xi_n) \neq 0, \\ 1 & \text{if } f(\xi_n) = 0, \end{cases} \quad 0 \le n < q.$$

They proved that if  $f(X) \in \mathbb{F}_q[X]$  has no multiple zero, and if either

(i) k = 2 and deg f < p

or

(ii)  $4^{r(k+\deg f)} < p$ ,

then

$$C_k(L_q) = O(2^{(r-1)k} r 2^r k \deg f q^{1/2} (\log p)^r),$$

By applying Corollary 1, we get in the same way as in [11] that the correlation measure satisfies

$$C_k(L_q) = O((k+1)^{r-1}k \deg fq^{1/2}(\log p)^r)$$

under the same conditions (i) and (ii).

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Moreover, if f(X) is a linear polynomial, then we get

$$C_k(L_q) = O((k+1)^{r-1}kq^{1/2}(\log p)^r)$$

without conditions on k of type (i) or (ii). Therefore by [2, Theorem 1] we get a non-trivial lower bound on the linear complexity profile of the sequence  $S_q = (s_0, s_1, \dots, s_{q-1}) \in \{0, 1\}^q$  defined by  $s_n = (l_n + 1)/2$ 

$$L(S_q, N) = \Omega\left(\frac{N^{1/r}}{p^{1/2}\log p}\right), \quad 2 \le N < q.$$

2.3. On the linear complexity profile of the threshold sequence for digital explicit inversive pseudorandom numbers

Let  $(\gamma_n)$  be a sequence of digital explicit inversive pseudorandom numbers defined by (4). If

$$\gamma_n = c_{n,1}\beta_1 + c_{n,2}\beta_2 + \cdots + c_{n,r}\beta_r$$

with  $0 \le c_{n,i} < p$  (for all *i*, *j*), we define *digital explicit inversive pseudorandom numbers* of period *q* in the interval [0, 1) by defining

$$y_n = \sum_{j=1}^{r} c_{n,j} p^{-j}, \quad n = 0, 1, \dots$$

We here derive a bound on the correlation measure of order k of the binary sequences  $E_q = (e_0, e_1, \dots, e_{q-1})$  defined by

$$e_n = \begin{cases} 1 & \text{if } 0 \le y_n < \frac{1}{2}, \\ -1 & \text{if } \frac{1}{2} \le y_n < 1, \end{cases} \quad n = 0, 1, \dots, q-1,$$

and also the linear complexity profile of the associated bit sequence  $R_q = (r_0, r_1, ..., r_{q-1}) \in \{0, 1\}^q$ , where  $r_n := (e_n+1)/2$ . Chen, Gomez, and Winterhof [3] studied the correlation measure of digital explicit inversive pseudorandom numbers and their linear complexity profile. They use the upper bound  $r2^r 2^{k(r-1)}$  on the number of possible boxes (3), and obtain that

$$C_k(E_q) = O\left(r2^r 2^{k(r-1)} kq^{1/2} (\log q)^k (1 + \log p)^r\right)$$

and

$$L(R_q, N) = \Omega\left(\frac{\log(Nq^{-1/2}2^{-r}r^{-1}(1+\log p)^{-r})}{r+\log\log q}\right), \quad 2 \le N < q.$$

Using Corollary 1 in the proof of [3] gives the improved bound:

$$C_k(E_q) = O\left((k+1)^r q^{1/2} (\log q)^k (1+\log p)^r\right)$$

and by [2, Theorem 1]:

$$L(R_q, N) = \Omega\left(\frac{\log(Nq^{-1/2}(1 + \log p)^{-r}(\log q^{1/2})^{-r}(\log \log q)^{r})}{\log \log q}\right)$$

for  $2 \leq N < q$ .

#### 3. Proofs

**Proof of Theorem 1.** We can assume that  $r \ge 2$  since otherwise the theorem is obvious. Let  $L_{i,1}$  and  $L_{i,2}$  be integers as follows:

$$L_{i,1} = \max\{0, pw_{1,i+1} - d_{1,i} - w_{1,i}, \dots, pw_{k,i+1} - d_{k,i} - w_{k,i}\}$$

and

$$L_{i,2} = \min\{p, pw_{1,i+1} - d_{1,i} - w_{1,i} + p, \dots, pw_{k,i+1} - d_{k,i} - w_{k,i} + p\}$$

for  $i \in \{1, 2, ..., r-1\}$ , where  $0 \le d_1 < \cdots < d_k < q$  and  $\omega_1, ..., \omega_k \in W$  such that  $d_j = \sum_{i=1}^r d_{j,i} p^{i-1}$ ,  $0 \le d_{j,i} < p$  and  $\omega_j = w_{j,2}\beta_2 + \cdots + w_{j,r}\beta_r$ . We also define  $L_{r,1} = 0$  and  $L_{r,2} = p$ .

Then clearly

$$S_{d_1,\omega_1} \cap \ldots \cap S_{d_k,\omega_k} = \{\ell_1\beta_1 + \cdots + \ell_r\beta_r : L_{i,1} \le \ell_i < L_{i,2}, i = 1, 2, \dots, r\}$$

is a box or empty.

In order to prove the bound on the number of nonempty sets having the form (2), we write (2) as an intersection of decreasing boxes, when the *k*-tuple  $(d_1, \ldots, d_k)$  is fixed.

First, let  $\overline{\omega}_t \in \{0, 1\}^k$  (for t = 1, ..., r) be a tuple of the *t*th coordinates of  $\omega_1, ..., \omega_k$ :  $\overline{\omega}_t = (w_{1,t}, ..., w_{k,t})$ . In particular,  $\overline{\omega}_1$  is the zero vector. For a  $\overline{\omega}_2 \in \{0, 1\}^k$  we define  $H^1(\overline{\omega}_2)$  as follows:

$$H^{1}(\overline{\omega}_{2}) = \{a \in \mathbb{F}_{p} : L_{1,1} \leq a < L_{1,2}\}.$$

It is clear that the sets in  $\{H^1(\overline{\omega}_2) : \overline{\omega}_2 \in \{0, 1\}^k\}$  split  $\mathbb{F}_p$ :

(i)  $H^1(\overline{\omega}_2) \subset \mathbb{F}_p$ ,

(ii) 
$$H^1(\overline{\omega}_2) \cap H^1(\overline{\omega}_2') = \emptyset$$
, for  $\overline{\omega}_2 \neq \overline{\omega}_2'$ 

(iii)  $\bigcup_{\overline{\omega}_2 \in \{0,1\}^k} H^1(\overline{\omega}_2) = \mathbb{F}_p.$ 

Moreover, there are at most k + 1 nonempty sets among them. Indeed, for an  $a \in \mathbb{F}_p$  we have

 $a > p - d_{i,1} \Leftrightarrow w_{i,2} = 1$ , for i = 1, 2, ..., k,

which means that there is a unique vector  $\overline{\omega}_2 \in \{0, 1\}^k$  such that  $a \in H^1(\overline{\omega}_2)$ . Let  $\sigma$  be a permutation of  $\{1, \ldots, k\}$  such that

 $d_{\sigma(1),1} \leq \cdots \leq d_{\sigma(k),1}$ 

Then, the possible sets of  $H^1(\overline{\omega}_2)$  are

(i)  $[0, p - d_{\sigma(k), 1}) \cap \mathbb{N}$ ,

(ii)  $[p - d_{\sigma(i),1}, p - d_{\sigma(i-1),1}) \cap \mathbb{N}$ , for i = 2, 3, ..., k,

(iii)  $[p - d_{\sigma(1),1}, p) \cap \mathbb{N}$ 

and the empty set.

Let

$$C^{1}(\overline{\omega}_{2}) = H^{1}(\overline{\omega}_{2}) \times \mathbb{F}_{p}^{r-1}$$
$$= \{\ell_{1}\beta_{1} + \ell_{2}\beta_{2} + \dots + \ell_{r}\beta_{r} : \ell_{1} \in H^{1}(\overline{\omega}_{2}), \ell_{2}, \dots, \ell_{r} \in \mathbb{F}_{p}\}.$$

By recursion we split a non-empty set  $C^{t-1}(\overline{\omega}_2, \ldots, \overline{\omega}_t)$   $(t = 2, 3, \ldots, r-1)$  as a disjoint union

$$C^{t-1}(\overline{\omega}_2,\ldots,\overline{\omega}_t) = \bigcup_{\overline{\omega}_{t+1} \in \{0,1\}^k} C^t(\overline{\omega}_2,\ldots,\overline{\omega}_t,\overline{\omega}_{t+1})$$

such that there are at most k + 1-many non-empty sets  $C^t(\overline{\omega}_2, \ldots, \overline{\omega}_t, \overline{\omega}_{t+1})$  when  $\overline{\omega}_{t+1}$  runs in  $\{0, 1\}^k$ .

Assume that  $C^{t-1}(\overline{\omega}_2,\ldots,\overline{\omega}_t)$  is defined for  $2 \le t < r-1$  and the k-tuples  $\overline{\omega}_2,\ldots,\overline{\omega}_t$  are fixed. Define the sets  $H^t(\overline{\omega}_{t+1})$  as

$$H^{t}(\overline{\omega}_{t+1}) = \{a \in \mathbb{F}_{p} : L_{t,1} \leq a < L_{t,2}\}.$$

The sets in  $\{H^t(\overline{\omega}_{t+1}) : \overline{\omega}_{t+1} \in \{0, 1\}^k\}$  split  $\mathbb{F}_p$ :

- (i)  $H^{t}(\overline{\omega}_{t+1}) \subset \mathbb{F}_{p}$ , (ii)  $H^{t}(\overline{\omega}_{t+1}) \cap H^{t}(\overline{\omega}'_{t+1}) = \emptyset$ , for  $\overline{\omega}_{t+1} \neq \overline{\omega}'_{t+1}$ ,
- (iii)  $\bigcup_{\overline{\omega}_{t+1} \in \{0,1\}^k} H^t(\overline{\omega}_{t+1}) = \mathbb{F}_p.$

Similar to the case t = 1, there are at most k + 1 nonempty sets among them. Indeed, for an integer  $a \in \mathbb{F}_p$  we have

$$a \ge p - d_{j,t} - w_{j,t} \Leftrightarrow w_{j,t+1} = 1, \quad \text{for } j = 1, \dots, k.$$

Then we define  $C^t$  ( $\overline{\omega}_2, \ldots, \overline{\omega}_t, \overline{\omega}_{t+1}$ ) as

$$\mathcal{C}^{t}(\overline{\omega}_{2},\ldots,\overline{\omega}_{t},\overline{\omega}_{t+1}) = \{\ell_{1}\beta_{1} + \ell_{2}\beta_{2} + \cdots + \ell_{r}\beta_{r} \in \mathcal{C}^{t-1}(\overline{\omega}_{2},\ldots,\overline{\omega}_{t}) : \ell_{t} \in \mathcal{H}^{t}(\overline{\omega}_{t+1})\}$$

We finish the proof of Theorem 1 by observing that

$$S_{d_1,\omega_1}\cap\ldots\cap S_{d_k,\omega_k}=C^{r-1}(\overline{\omega}_2,\ldots,\overline{\omega}_r).$$

**Proof of Corollary 1.** Let  $M \in \{0, 1, ..., q - 1\}$  and  $d_1, ..., d_k < q$  with  $0 \le d_1 < \cdots < d_k < q$  be given integers. We look for the intersection of  $S_{d_1,\omega_1} \cap \cdots \cap S_{d_k,\omega_k}$  and the set  $\mathcal{E} = \{\xi_0, \ldots, \xi_{M-1}\} \subseteq \mathbb{F}_q$  for tuples  $(\omega_1, \ldots, \omega_k) \in \mathcal{W}^k$ . This intersection has at most r boxes for only the tuple  $(\omega_1^0, \ldots, \omega_k^0) \in \mathcal{W}^k$  where  $\xi_{M-1} \in S_{d_1,\omega_1^0} \cap \cdots \cap S_{d_k,\omega_k^0}$ . Moreover, there are (k + 1) tuples for which the intersection consists of at most (r - 1) boxes. By Theorem 1 we know that we have at most  $(k + 1)^{r-1}$  tuples for which  $S_{d_1,\omega_1} \cap \cdots \cap S_{d_k,\omega_k}$  is nonempty. Therefore, in the final step there are  $(k + 1)^{r-1}$  tuples whose intersection with  $\{\xi_0, \xi_2, \dots, \xi_{M-1}\}$  consists of at most one box. If we sum them up we obtain that

$$\sum_{i=0}^{r-1} (r-i)(k+1)^i = \frac{(k+1)^{r+1} - (r+1)k - 1}{k^2} = O((k+1)^{r-1}).$$

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